

Constraint Logic Programming over Infinite Domains with an Application to Proof

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Background



- Model Checking
- Data Validation
- Proof and Disproof

B Example

MACHINE BinarySearch

CONSTANTS arr, size, goal

PROPERTIES

size : NAT & arr : 1..size \longrightarrow INT &
!(i).(i:1..(size-1) \Rightarrow arr(i) \leq arr(i+1)) &
goal : INT &
arr = [-1,0,1,2,3,5] & goal : {5,4,2,6}

VARIABLES i, j, found

INVARIANT

found:BOOL & i:NAT & j:NAT &
(found=TRUE \Rightarrow arr(i)=goal) &

INITIALISATION found, i, j := FALSE, 1, size

OPERATIONS

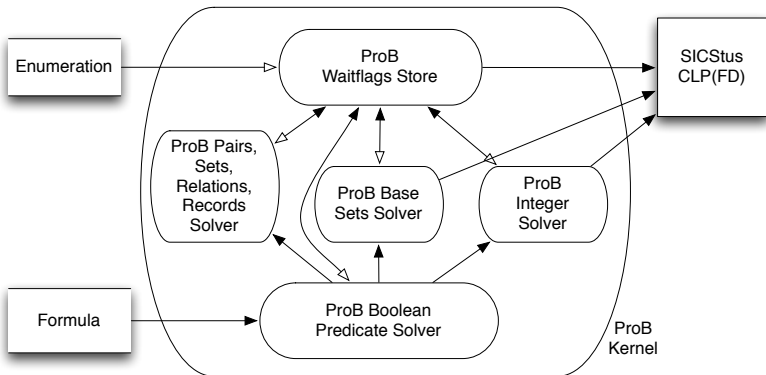
Step = SELECT found=FALSE & i<j THEN

...

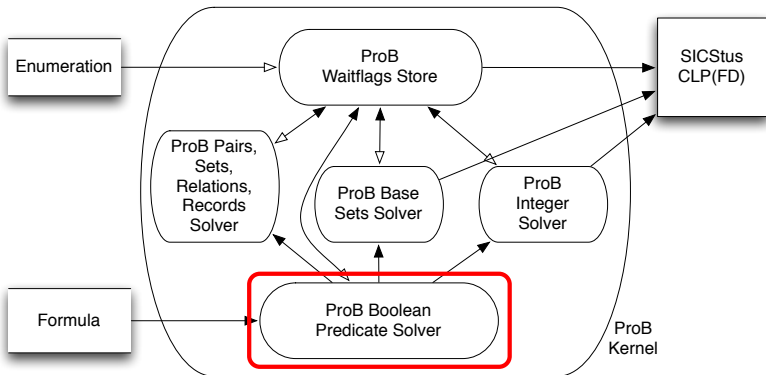
PROB Kernel

- Based on SICStus Prolog
 - 150k LoC
 - + Extensions in Tcl/Tk, Java, C, ...
- CLP(FD) + CHR
- Custom solvers for high-level structures
 - Sets, lists, tuples, records, ...
 - Nested quantification
 - Comprehensions & lambdas

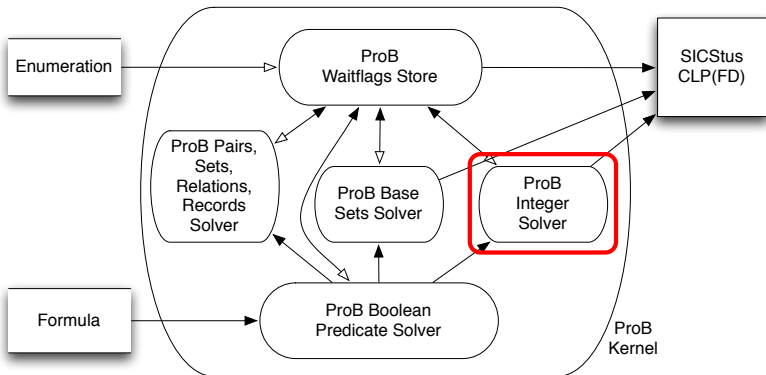
ProB Kernel



ProB Kernel



ProB Kernel



Techniques for Large or Infinite Domains

- Tracking of enumeration scopes
- Randomized enumeration
- High-level reasoning

Types of Enumeration

Distinguish based on effect on overall result:

- No enumeration, result not influenced
 - e.g., no valuation \Rightarrow formula unsatisfiable
- Exhaustive enumeration
- Non-exhaustive enumeration with or w/o result
 - \Rightarrow consider scope

Nested Enumeration Scopes

$\exists x.(x \in [-10, 10] \wedge \forall y.y \in [-15, 5] \Rightarrow y \leq x)$							
existential scope							
x can be enumerated exhaustively but is not							
	<table border="1"> <tr> <td>$\forall y.y \in [-15, 5] \Rightarrow y \leq x$</td> </tr> <tr> <td>universal scope</td> </tr> <tr> <td>y is enumerated exhaustively</td> </tr> <tr> <td> <table border="1"> <tr> <td>$y \leq x$</td> </tr> <tr> <td>CLP(FD) constraint</td> </tr> </table> </td> </tr> </table>	$\forall y.y \in [-15, 5] \Rightarrow y \leq x$	universal scope	y is enumerated exhaustively	<table border="1"> <tr> <td>$y \leq x$</td> </tr> <tr> <td>CLP(FD) constraint</td> </tr> </table>	$y \leq x$	CLP(FD) constraint
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CLP(FD) constraint							

Nested Enumeration Scopes

$\exists x. (\neg \forall y. y \in [-15, 5] \Rightarrow y \leq x)$			
existential scope			
non-exhaustively search for x			
$\forall y. y \in [-15, 5] \Rightarrow y \leq x$			
existential scope			
y can be enumerated exhaustively but is not			
	<table border="1"> <tr> <td style="text-align: center;">$y \leq x$</td> </tr> <tr> <td>CLP(FD) constraint</td> </tr> </table>	$y \leq x$	CLP(FD) constraint
$y \leq x$			
CLP(FD) constraint			

Examples

- $x * x = 10000$, solution found w/o enumeration
- $\{x \mid x * x = 10000\}$, (all) solutions found w/o enumeration
- $x > 10000 \ \& \ x \bmod 1234 = 1$, solution found despite enumeration
- $\{x \mid x > 10000 \ \& \ x \bmod 1234 = 1\}$ can thus not be solved
- $x * x = 10001$ unsat by CLP(FD) propagation w/o enumeration
- $x > 10000 \ \& \ x \bmod 1234 = 1 \ \& \ x * x = 10 * x$ partial enumeration, no result

Randomized Enumeration

- Try values of different size / characteristics
- Less likely to get stuck in some part of the search space
- Increase variance, e.g., for test case generation

Randomized Enumeration - Challenges

- Random permutation
 - \Rightarrow avoid duplicates
- Detect exhaustive traversal

Fisher-Yates / Knuth Shuffle!

Data: List a

Result: Random permutation of a

procedure shuffle(a)

for $i \in [0, \text{length}(a) - 1]$ **do**

 chose j randomly such that $0 \leq j \leq i$

if $j \neq i$ **then**

$\text{perm}[i] := \text{perm}[j]$

end

$\text{perm}[j] := a[i]$

end

return perm

Fisher-Yates / Knuth Shuffle :(

Data: List a

Result: Random permutation of a

procedure shuffle(a)

for $i \in [0, \text{length}(a) - 1]$ **do**

 chose j randomly such that $0 \leq j \leq i$

if $j \neq i$ **then**

$\text{perm}[i] := \text{perm}[j]$

end

$\text{perm}[j] := a[i]$

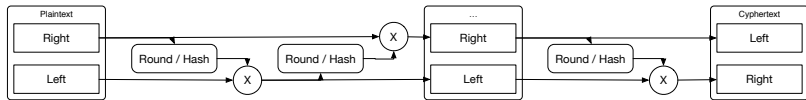
end

return perm

Randomized Enumeration - Solution

- Use techniques from cryptography!
- “Encrypt” $[l, u] \rightarrow [l, u]$
- Using random encryption key
 - \Rightarrow random enumeration of $[l, u]$
- Unique encryption
 - \Rightarrow no duplicates, random permutation

Feistel Network



High-Level Reasoning using CHR

- CLP(FD) weak without bounds
- e.g. no unsat for $X < Y \wedge Y < X$
- High-level reasoning:
 - Infer falsity
 - Find new CLP(FD) constraints

CHR Rules (excerpt)

reflexivity @ $\text{leq}(X,X) \Leftrightarrow \text{true}$.

antisymmetry @ $\text{leq}(X,Y), \text{leq}(Y,X) \Leftrightarrow X = Y$.

idempotence @ $\text{leq}(X,Y) \wedge \text{leq}(X,Y) \Leftrightarrow \text{true}$.

transitivity @ $\text{leq}(X,Y), \text{leq}(Y,Z) \Rightarrow \text{leq}(X,Z)$.

antireflexivity @ $\text{lt}(X,X) \Leftrightarrow \text{fail}$.

idempotence @ $\text{lt}(X,Y) \wedge \text{lt}(X,Y) \Leftrightarrow \text{true}$.

transitivity @ $\text{lt}(X,Y), \text{leq}(Y,Z) \Rightarrow \text{lt}(X,Z)$.

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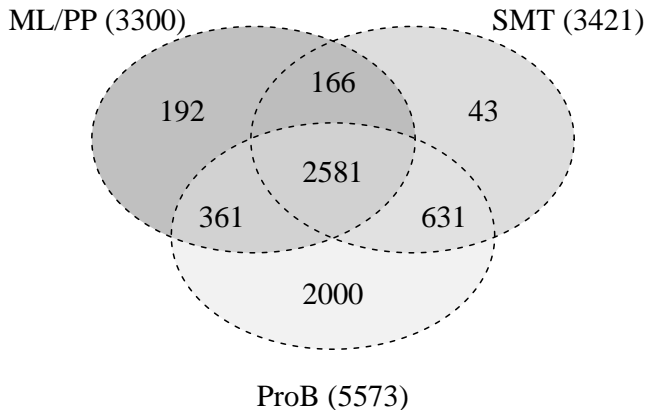
Comparison

predicate	with CHR	without CHR
$x > 3$	1.2	1.08
$x > y \wedge y > x$	0.92	timeout
$x = 3 \wedge x > y \wedge y = 4$	0.98	1.07
$x = 3 \wedge x > y$	0.86	1.0
$x = 3 \wedge x < y$	1.1	1.11
$w > x \wedge x > y \wedge y > z \wedge w = 1 \wedge z = 1$	0.98	0.95
$w > x \wedge x > y \wedge y > z \wedge z > w$	0.88	timeout
$x + 2 > y + 1 \wedge y > x$	timeout	timeout
$x > y \wedge y > x + 1$	0.9	timeout

Application - Proof

- Given Event-B sequent $H_i(x_1, \dots, x_k) \models G(x_1, \dots, x_k)$
- Build $\exists x_1, \dots, x_k : (H_1(x_1, \dots, x_k) \wedge \dots \wedge H_n(x_1, \dots, x_k)) \Rightarrow \neg G(x_1, \dots, x_k)$
- Search for counter-example
- Exhaustive search \Rightarrow proof

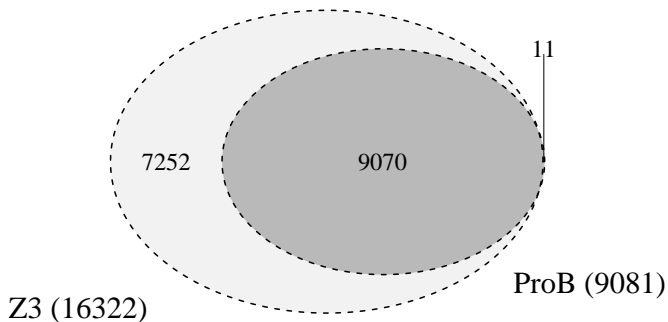
Results



Application - SMT Solving

- SMT-LIB standard fully supported
- Translation from SMT-LIB to B
- Competitive on integer arithmetic
- Not competitive for bitvectors, arrays and sets

Results on Integer Arithmetic



SMT-COMP 2016 - NIA - Main Track

Winners:

Sequential Performances	Parallel Performances
ProB	ProB

Result table¹

Solver	Sequential performance		
	Error Score	Correctly Solved Score	avg. CPU time
CVC4	0.000	5.000	0.019
ProB	0.000	8.000	0.888
vampire_smt_4.1	0.000	6.000	335.043
vampire_smt_4.1_parallel	0.000	7.000	537.677
z3 [¶]	0.000	9.000	0.038

Thank you for your attention!
Are there any questions left?

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