

Termination of $\{aa \rightarrow cb, bb \rightarrow ca, cc \rightarrow ba\}$

Dieter Hofbauer (Kassel) and Johannes Waldmann (Leipzig)

3 August 05

The general setting: Let \succ be a well-founded order on a non-empty set D , and let \succeq denote its reflexive closure. Let $i : \Sigma \rightarrow (D \rightarrow D)$ be a mapping which is extended to Σ^* by $i(\epsilon)(d) = d$ and $i(xs)(d) = i(x)(i(s)(d))$ for $x \in \Sigma$, $s \in \Sigma^*$, $d \in D$. Further assume strict monotonicity, i.e., for $x \in \Sigma$ and $d, d' \in D$, $d \succ d'$ implies $i(x)(d) \succ i(x)(d')$. For a proof of the following standard result see e.g. Theorem 4 in [Zantema, Termination of string rewriting proved automatically].

Lemma 1. *Let R and S be string rewriting systems over alphabet Σ such that*

- $i(\ell)(d) \succ i(r)(d)$ for all $\ell \rightarrow r$ in R and $d \in D$,
- $i(\ell)(d) \succeq i(r)(d)$ for all $\ell \rightarrow r$ in S and $d \in D$.

Then R is terminating relative to S .

Proposition 1. $Z_{086} = \{aa \rightarrow cb, bb \rightarrow ca, cc \rightarrow ba\}$ is terminating.

Proof. Let $\Sigma = \{a, b, c\}$. Consider the interpretation $i : \Sigma \rightarrow (\mathbb{N}^4 \rightarrow \mathbb{N}^4)$ with

$$\begin{aligned} i(a)(u, v, x, y) &= (u + y, y, v + 2y + 2, v), \\ i(b)(u, v, x, y) &= (u + v, 2v + y + 1, v, 0), \\ i(c)(u, v, x, y) &= (u + 2y, x + y, v + 2y + 1, 0), \end{aligned}$$

Let $(u, v, x, y) \succ (u', v', x', y')$ if $u > u'$ and $v \geq v'$, $x \geq x'$, $y \geq y'$, and let \succeq be the reflexive closure of \succ . Note that \succ is well-founded on \mathbb{N}^4 , and that $i(x)$ is strictly monotone for each $x \in \Sigma$. As is easily verified,

$$\begin{aligned} i(aa)(u, v, x, y) &= (u + v + y, v, 2v + y + 2, y) \\ &\succeq (u + v, v, 2v + y + 2, 0) = i(cb)(u, v, x, y), \\ i(bb)(u, v, x, y) &= (u + 3v + y + 1, 4v + 2y + 3, 2v + y + 1, 0) \\ &\succ (u + 2v + y, 2v + 2y + 2, 2v + y + 1, 0) = i(ca)(u, v, x, y), \\ i(cc)(u, v, x, y) &= (u + 2y, v + 2y + 1, x + y + 1, 0) \\ &\succeq (u + 2y, v + 2y + 1, y, 0) = i(ba)(u, v, x, y). \end{aligned}$$

By the lemma above, termination of $S = \{aa \rightarrow cb, cc \rightarrow ba\}$ implies termination of Z_{086} . Termination of S is trivial: choose weights $a, c \mapsto 1$ and $b \mapsto 0$. \square

To illustrate, here ist the sequence of interpretations for a rewrite sequence starting from $aabb$:

$$\begin{array}{cccc} a(4, 3, 8, 0) & c(4, 3, 8, 0) & c(2, 3, 5, 0) & b(2, 3, 1, 0) \\ a(1, 0, 5, 3) & b(4, 7, 3, 0) & c(2, 4, 3, 0) & a(1, 1, 4, 0) \\ b(1, 3, 0, 0) \rightarrow_1 b(1, 3, 1, 0) \rightarrow_2 a(0, 0, 3, 1) \rightarrow_3 a(0, 0, 3, 1) \\ b(0, 1, 0, 0) & b(0, 1, 0, 0) & b(0, 1, 0, 0) & b(0, 1, 0, 0) \\ (0, 0, 0, 0) & (0, 0, 0, 0) & (0, 0, 0, 0) & (0, 0, 0, 0) \\ & b(2, 3, 1, 0) & b(1, 3, 1, 0) & b(1, 3, 1, 0) \\ & c(1, 1, 4, 0) & c(0, 1, 3, 0) & b(0, 1, 0, 0) \\ \rightarrow_4 b(1, 3, 1, 0) \rightarrow_5 c(0, 2, 1, 0) \rightarrow_6 a(0, 0, 2, 0) \\ & b(0, 1, 0, 0) & a(0, 0, 2, 0) & a(0, 0, 2, 0) \\ & (0, 0, 0, 0) & (0, 0, 0, 0) & (0, 0, 0, 0) \end{array}$$

Note the strict decreases in steps 2 and 5.